Async Rust and 9p server

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Agenda

- Introduction
  - Motivation, async Rust, kernel implementations
- 9p server
  - Description, some implementation samples
- Demo
- Discussion

Motivation

- Showcase productivity gains from using Rust
  - In addition to security benefits
- Considerable attack surface
  - For example, receiving untrusted data over the network
- Pure software module
  - No inherent unsafety due to bus mastering devices
- Started looking at `ksmbd`
  - Not an ideal initial project because of complexity of the protocol
  - Also requires a user-space component

Async Rust

Talked about it at OSS North America: link

In summary:

- Compiler automatically creates a state machine from thread-like code
- Kernel crate implements executors and reactors

Workqueue Executor

Spawning tasks

- Allocates task: contains future plus executor-specific state (e.g., `work_struct`)
- Adds to task list
- Wakes task up

Waking tasks up

- Enqueues task for running (e.g., `queue_work_on`)
- On worker thread: accesses revocable task, poll future, cleans it up when it completes

Tearing down

- All state is dropped (more on this later)

Socket Reactor

- Initialisation
  - Pinned larger struct containing some state plus wait queue entry (`wait_queue_entry`)
  - Wait queue entry with a custom function (`init_waitqueue_func_entry`)
  - Adds entry to the socket's wait queue (`add_wait_queue`)
- Waking up
  - Wait queue callback is called: uses `container_of` to get to outer struct
  - Checks mask for filter callbacks (`EPOLLIN`, `EPOLLOUT`, etc)
  - Calls `Waker::wake` to instruct executor to run task again
- Cleaning up
  - Removes entry from socket's wait queue (`remove_wait_queue`)

9p file server

What is 9p?

Plan 9 is an operating system from Bell Labs (link)

- Originally designed by Ken Thompson, Rob Pike, et al.

Included a remote file system protocol: Plan 9 File Protocol, 9P

- The Linux kernel already includes a 9p client
- Qemu implements a server to share a host directory with guest

Straightforward: original protocol only includes 10 operations

What is implemented

Not a file system

Though we have some support for it here

Exposes the local file system over the network

Read-only for now

Implements `9P2000.L` – Linux extensions

WIP but available here

Code stats

```
fs/k9pd/Kconfig     |  10 +
fs/k9pd/Makefile    |   5 +
fs/k9pd/buffer.rs   | 286 ++++++++++++++++
fs/k9pd/k9pd.rs     | 171 +++++++++
fs/k9pd/localfs.rs  | 382 ++++++++++++++++++++++
fs/k9pd/protocol.rs | 200 ++++++++++
6 files changed, 1054 insertions(+)
```

Receiving requests

```
async fn next_pdu(&self, max_size: u32) -> Result<Box<[u8]>> {
    // Read the length.
    let mut len_in_bytes = [0u8; 4];
    self.stream.read_all(&mut len_in_bytes).await?;

    let len = u32::from_le_bytes(len_in_bytes).checked_sub(4).ok_or(EIO)?;
    if len > max_size {
        return Err(E2BIG);
    }

    // Allocate the buffer and read the rest.
    self.stream.alloc_read_exact(len as usize).await
}
```

Dispatching requests

```
loop {
    let pdu = conn.next_pdu(max_size).await?;
    let (_op, tag) = protocol::get_op_tag(&pdu)?;
    let res = spawn_task!(
        executor.as_ref_borrow(), conn.clone().handle_pdu(tag, pdu));
    if let Err(e) = res {
        conn.write_result(|b| protocol::error(b, tag, e)).await?;
    }
}
```

Serialising writes

```
async fn write(&self, buf: &[u8]) {
    let mut inner = self.inner.lock().await;
    if inner.err.is_err() {
        // A previous write failed so we won't even try this one.
        return;
    }
    let ret = self.stream.write_all(buf).await;
    if ret.is_err() {
        // Store away the error.
        inner.err = ret;
    }
}
```

Cleaning up

Executor keeps track of all incomplete tasks

Auto-stop handles stop executors when they go out of scope

Stopping an executor waits for running tasks to go to sleep
- Also ensures that sleeping tasks don't wake up anymore
- Drops all tasks

"Local variables" of async functions are also dropped
- Reactors unregister from subsystems
- Allocations are freed, ref-counted objects are released, etc.

Stopping an executor results in everything being dropped automatically

Demo

Discussion

Possible topics

- Additional potential candidates to be written in async Rust
  - Driver state machines?
- Additional reactors
  - kiocb, urb, bio?
- Implementing more non-blocking operations/primitives
  - Select
  - Communication channels (send/receive data structures)
  - Memory allocation (GFP_KERNEL): can we give up the worker instead of sleeping?
  - Read-write mutexes
  - Condition variables

Possible topics (cont'd)

- Integration with user-space implementations
  - Tokio panics on memory allocation failures
- Destruction of mutual exclusion primitives
  - State may be inconsistent
  - Release the lock? Keep it locked? Poising (which implies acquisition may fail)?
- Fault isolation
  - Should we unwind stacks, drop objects, etc. for panics in tasks?
  - At the moment, we kill a workqueue thread

Thank you!